

Crunch Budapest October 2022

Apache Iceberg Merge-on-Read Streaming CDC

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AGENDA

- 1 Lakehouse Team's Mission**
- 2 Legacy System**
- 3 Our Design**
- 4 Result**
- 5 Reflection**

Lakehouse Team's Mission

*“Provide Shopify developers with an interoperable,
performant, and standards-based data lakehouse
where 1st- and 3rd-party Shopify data can be ingested”*

Table Snapshot

The state of a datasource at a specific moment in time.

- One row per a primary key
- Latest version of a primary key

Legacy System

Batch data ingestion
via statement based replication

Legacy System

“statement based replication”

→ Query DB for data

- ◆ Does not scale as tables get larger
- ◆ Long running queries (volatile)

```
SELECT * FROM my_table
```

Legacy System

“statement based replication”

- Query DB for data
 - ◆ Does not scale as tables get larger
 - ◆ Long running queries (volatile)
- Done incrementally by keeping position
- Depends on app devs updating updated_at
 - ◆ Possible to miss updates

```
SELECT * FROM my_table
```

```
WHERE updated_at > 2020/10/01 09:00:00
```

Legacy System

“statement based replication”

- Query DB for data
 - ◆ Does not scale as tables get larger
 - ◆ Long running queries (volatile)
- Done incrementally by keeping position
- Depends on app devs updating updated_at
 - ◆ Possible to miss updates
- Smaller queries using bucketing
 - ◆ Increased network requests increase time
- Can't capture deletes

```
SELECT * FROM my_table
```

```
WHERE updated_at > 2020/10/01 09:00:00
```

```
AND primary_key >= last_seen_key
```

```
ORDER BY primary_key
```

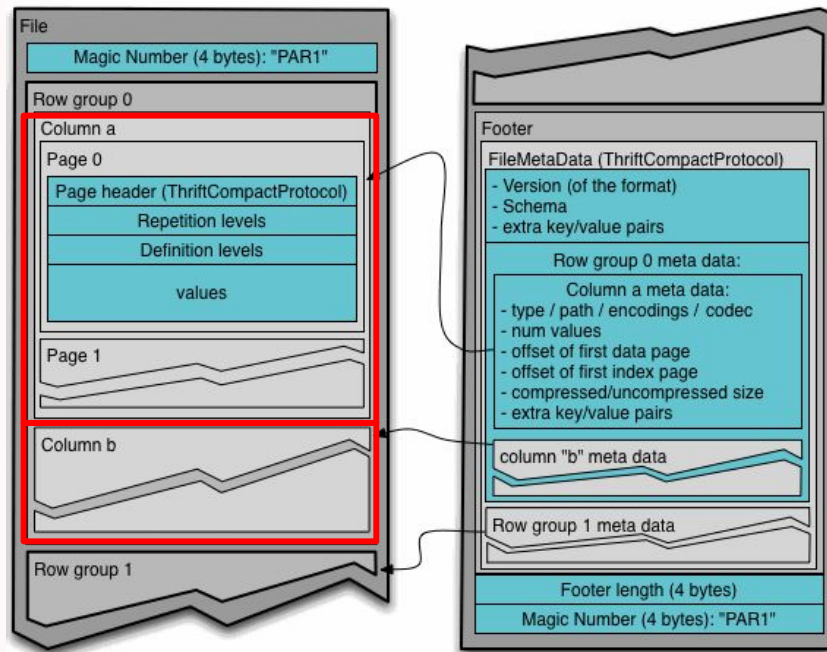
```
LIMIT 10,000
```


Legacy System

“statement based replication” contd.

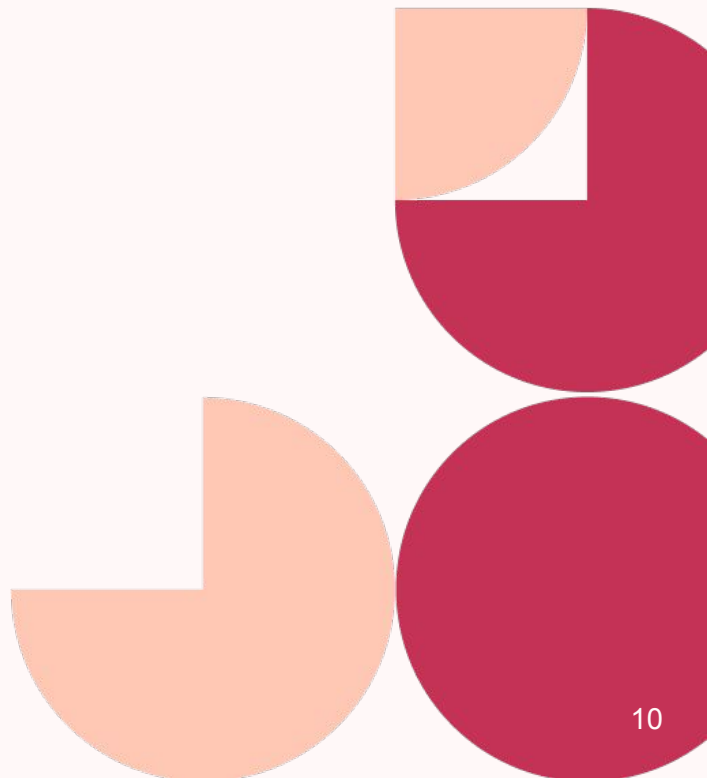
→ Spark process to apply update (⌚💰)

- ◆ Due to columnar file format
- ◆ Optimizing for aggregation analytics over a subset of columns
- ◆ Efficient compaction (schematized data)
- ◆ Columnar files are immutable (overwrite)
 - Rewrite is an expensive operation



Design Goals

- ❏ Scalable ingestion
- ❏ Scalable **snapshot production**
- ❏ More accurate updates
- ❏ Delete capture
- ❏ Snapshot SLO of under an hour

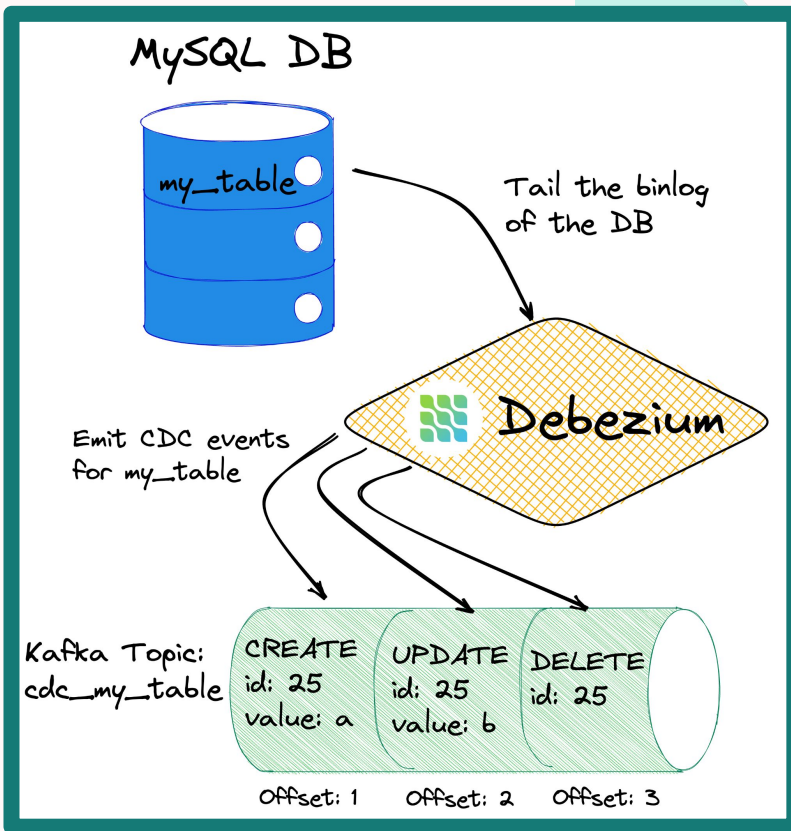


Design

*Streaming data ingestion of change data
capture via Kafka written in Iceberg V2 storage
format*

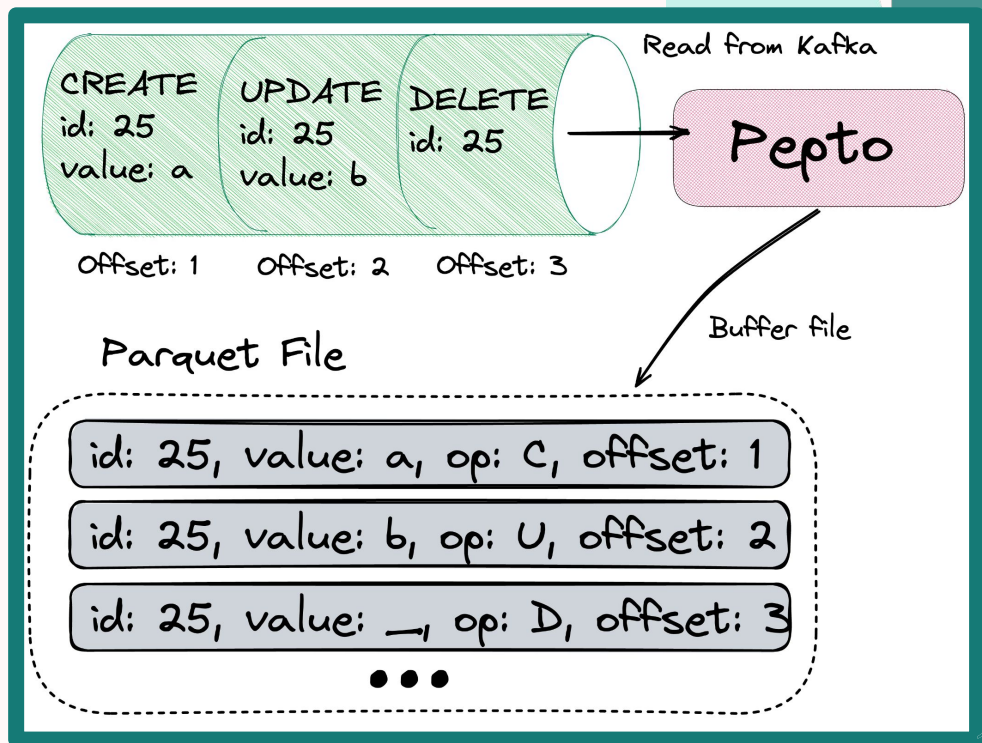
Design : Kafka ingestion + CDC

- Capturing changes from our source
 - Binlog from MySQL captures every transaction
 - **Every CREATE, UPDATE, DELETE**
 - Requires performing an upsert
 - Ordered based on when the event happened
 - Binlog is also used for DB replication
- State changes are emitted to Kafka for future ingestion



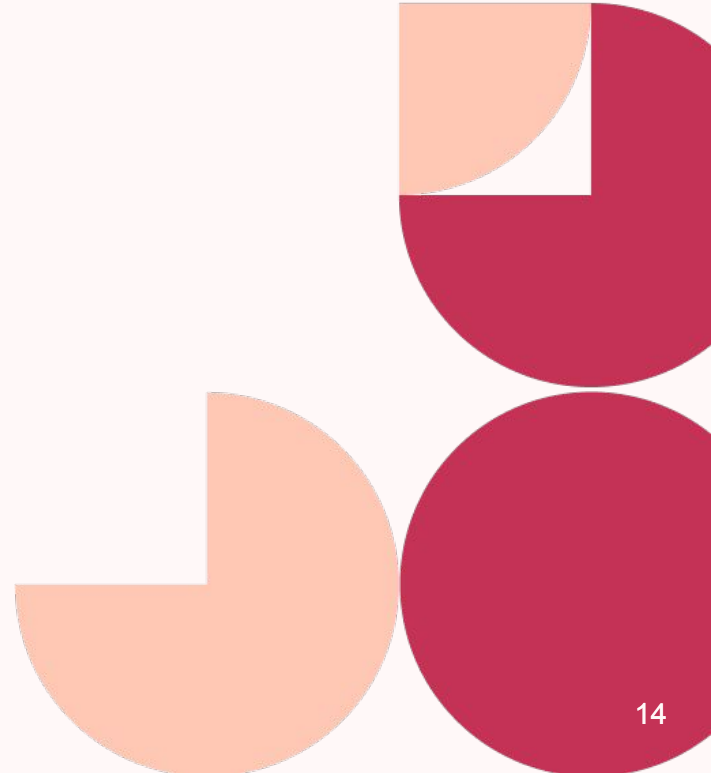
Design : Kafka ingestion + CDC

- Consume the CDC events from Kafka
- Buffer events into Parquet files
 - Registered in an Iceberg Table
- Achievements
 - Scalability w/ Kafka partitions + multiple consumers / writers
 - More accurate updates
 - Delete capture
 - ~7 min SLO data ingest



Goals

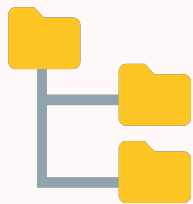
- ☒ Scalable ingestion
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Design : Iceberg

- Iceberg is a table format
- Just a library
- Contents of a table are identified by traversing through metadata files

File system



Hard Drive



Iceberg

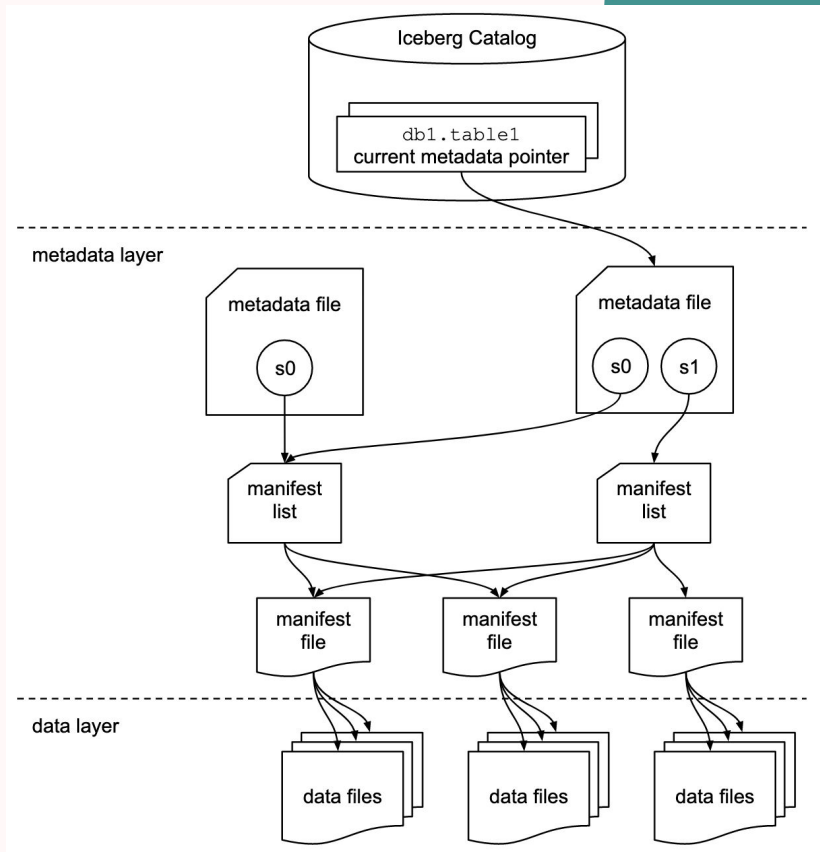


Object Storage



Design : Iceberg

- Iceberg is a table format
- Just a library
- Contents of a table are identified by traversing through metadata files



Design : Iceberg

- Iceberg is **metadata rich table format**
on top of parquet files
- Enables more **efficient file pruning**
 - More scalable reads
 - More scalable writes

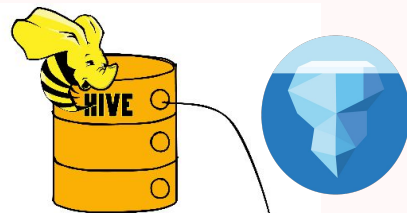


efficient read & writes

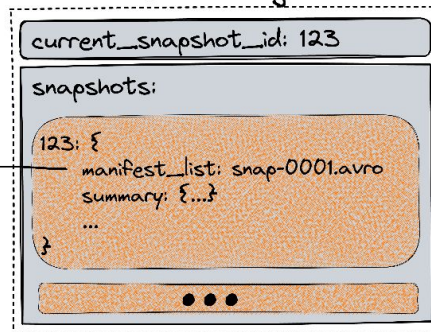


scalable snapshot production & 1h< SLO

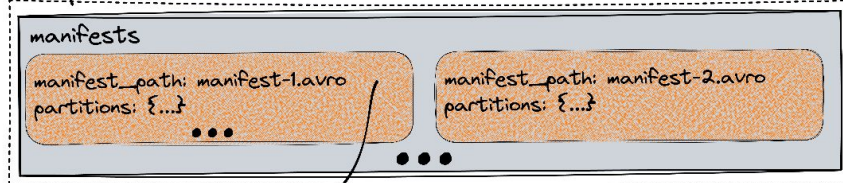
Design : Iceberg



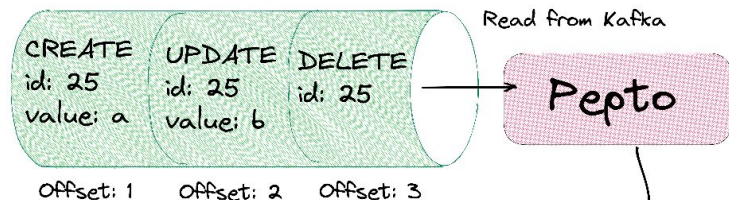
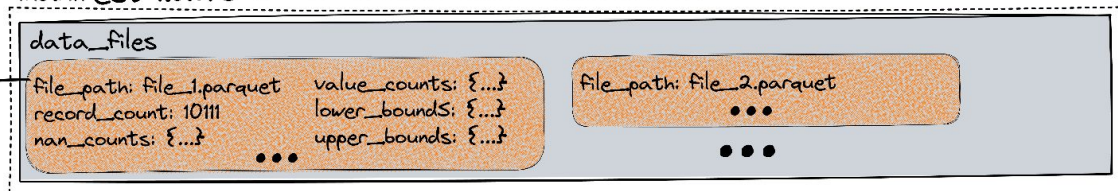
0001.metadata.json



snap-0001.avro

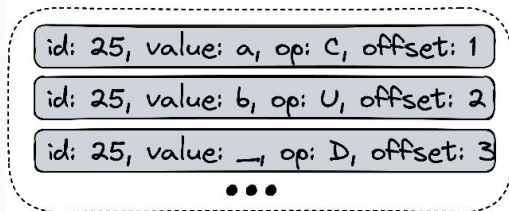


manifest-1.avro

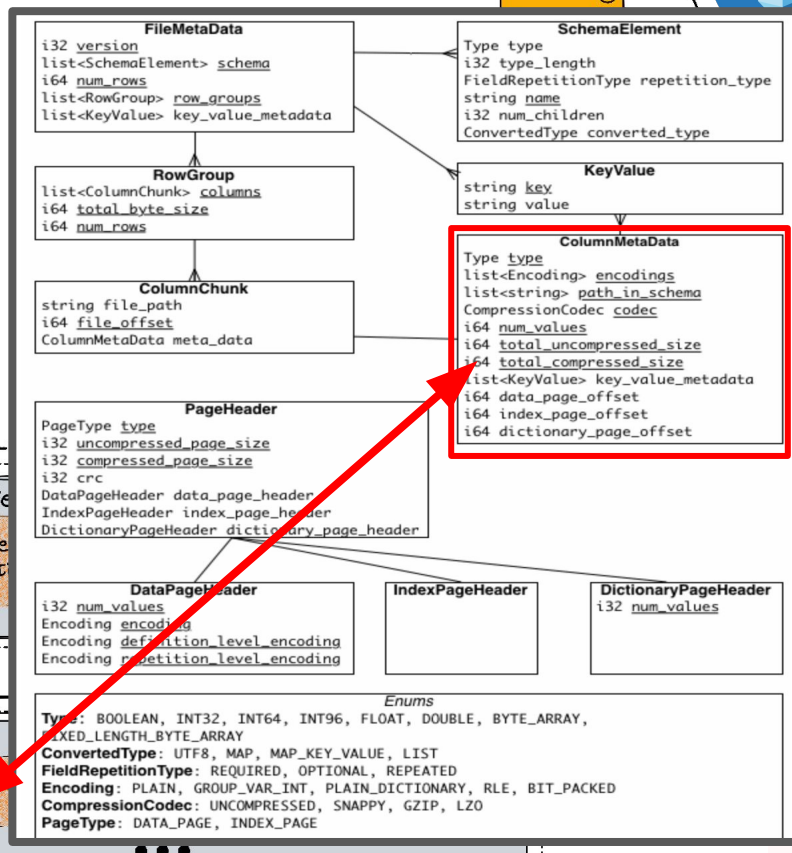
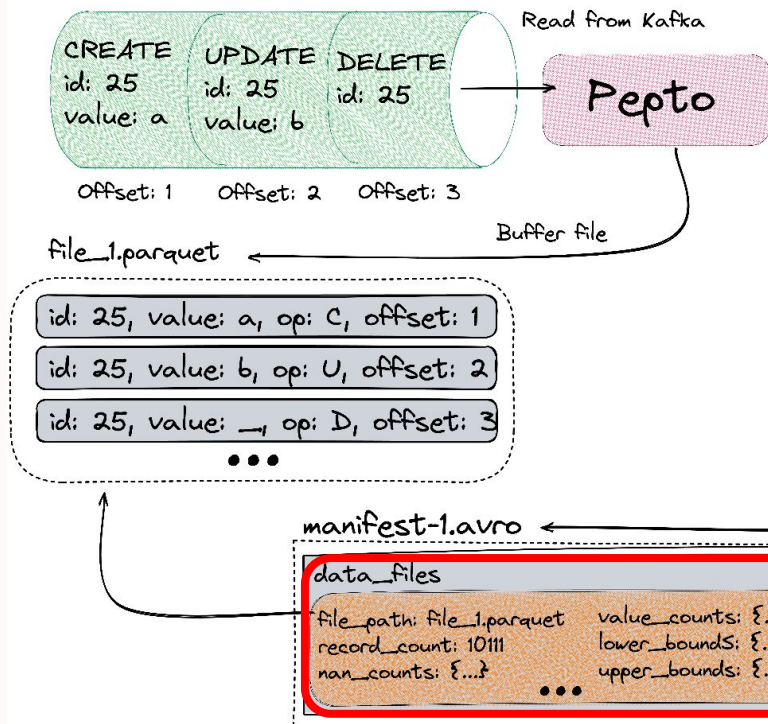


Buffer file

file_1.parquet



Design : Iceberg



Design

Kafka Injection

+

CDC

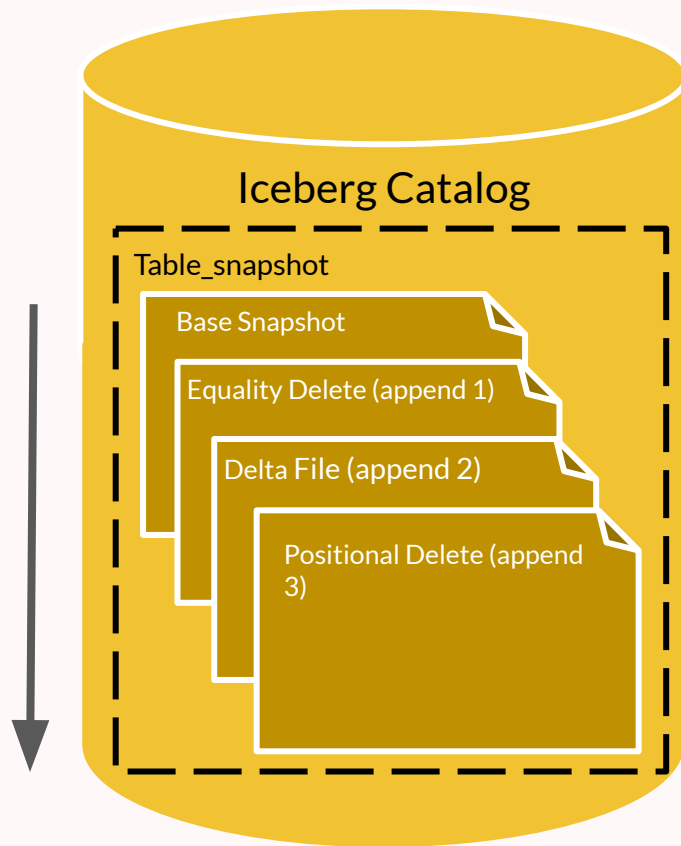
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Merge-on-Read Iceberg

(aka Iceberg V2 storage)

Design : Iceberg V2 storage

- V2 Spec introduces delete files
 - Positional Delete
 - Equality Delete
- Act as filters at query time
- **Procrastinate rewriting files**
- Targeted rewrite (via Iceberg metadata)
 - **Reduced compute**



Design : Iceberg V2 storage

- **Procrastinate** rewriting files
- Targeted rewrite (via Iceberg metadata)
 - **Reduced compute**



Targeted upsertes



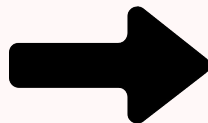
scalable snapshot production & 1h< SLO

Design : Iceberg V2 storage – Our Goal

Base Table		
id: 25	value: a	op: C
id: 30	value: alpha	op: U

+

Data File (file_1.parquet)		
id: 25	value: b	op: U
id: 45	value: c	op: C
id: 45	value: d	op: U



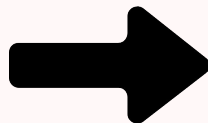
Desired Result		
id: 25	value: b	op: U
id: 30	value: alpha	op: U
id: 45	value: d	op: U

Design : Iceberg V2 storage – Our Goal

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Data File (file_1.parquet)		
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id: 45	value: d	op: U



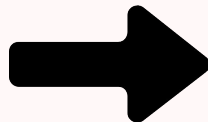
Desired Result		
id: 25	value: b	op: U
id: 30	value: alpha	op: U
id: 45	value: d	op: U

Design : Iceberg V2 storage – Step 1: Append Equality Delete

Base Table		
id: 25	value: a	op: C
id: 30	value: alpha	op: U

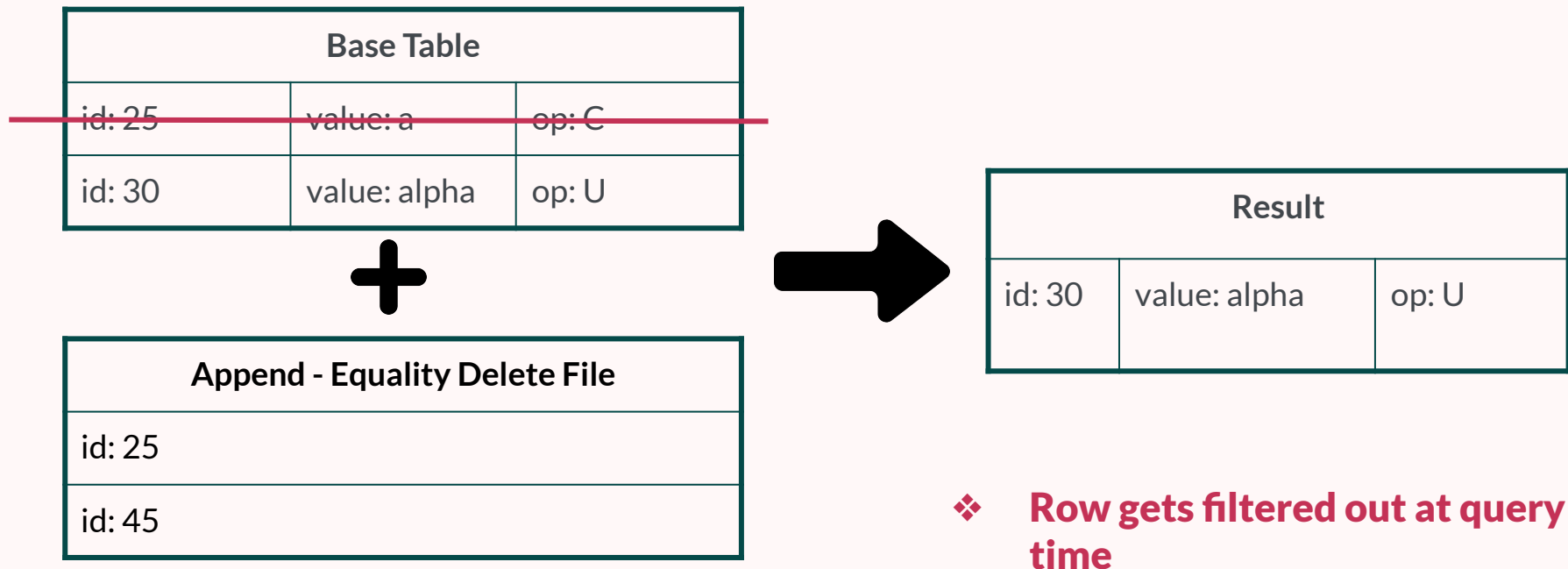
+

Append - Equality Delete File		
id: 25		
id: 45		

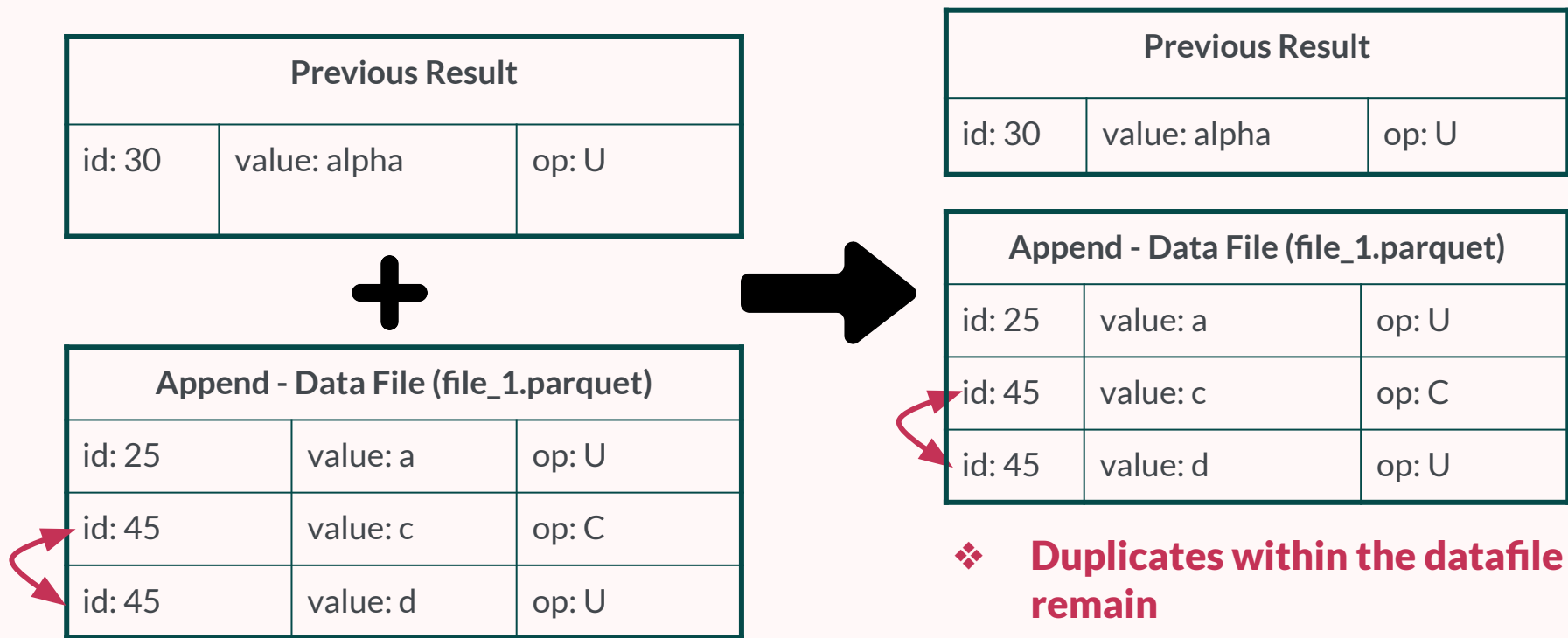


Result		
id: 30	value: alpha	op: U

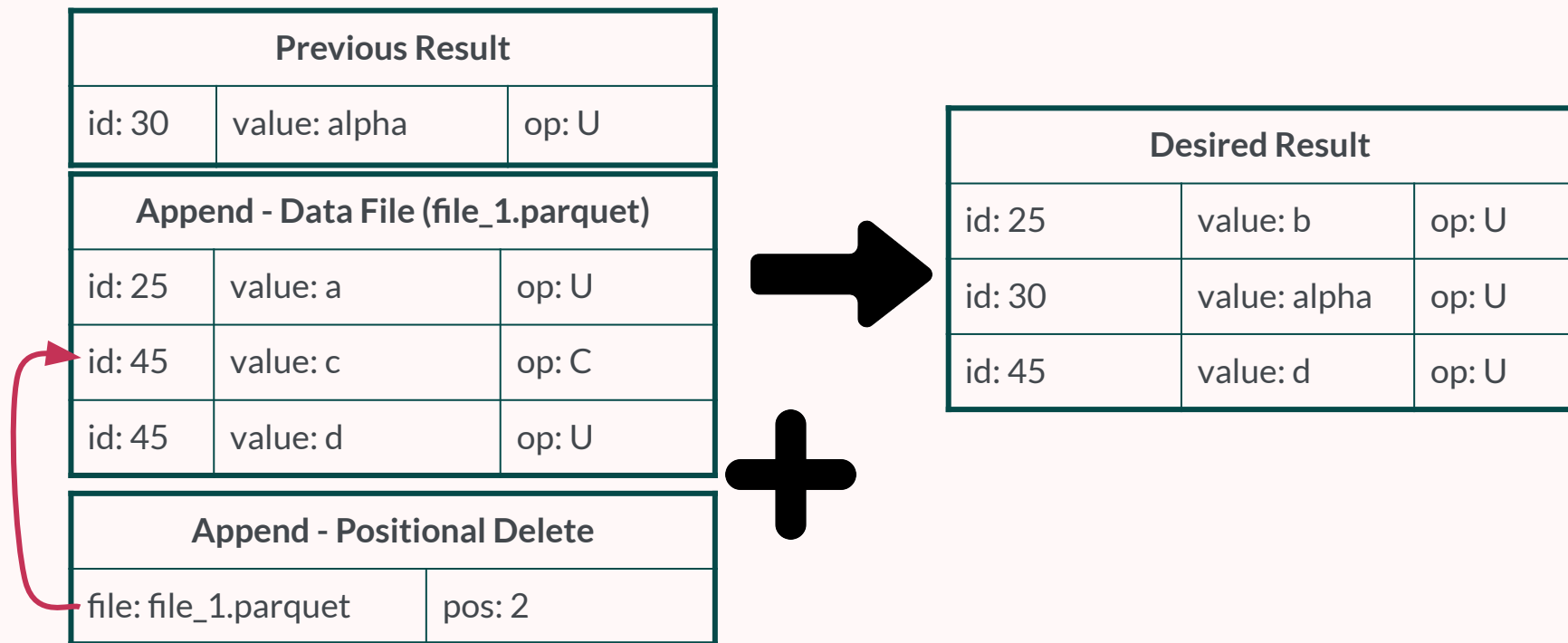
Design : Iceberg V2 storage – Step 1: Append Equality Delete



Design : Iceberg V2 storage – Step 2: Append Data File



Design : Iceberg V2 storage – Step 3: Append Positional Delete

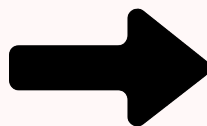


Design : Iceberg V2 storage – Step 3: Append Positional Delete

Previous Result		
id: 30	value: alpha	op: U

Append - Data File (file_1.parquet)		
id: 25	value: a	op: U
id: 45	value: c	op: C
id: 45	value: d	op: U

Append - Positional Delete		
file: file_1.parquet		pos: 2

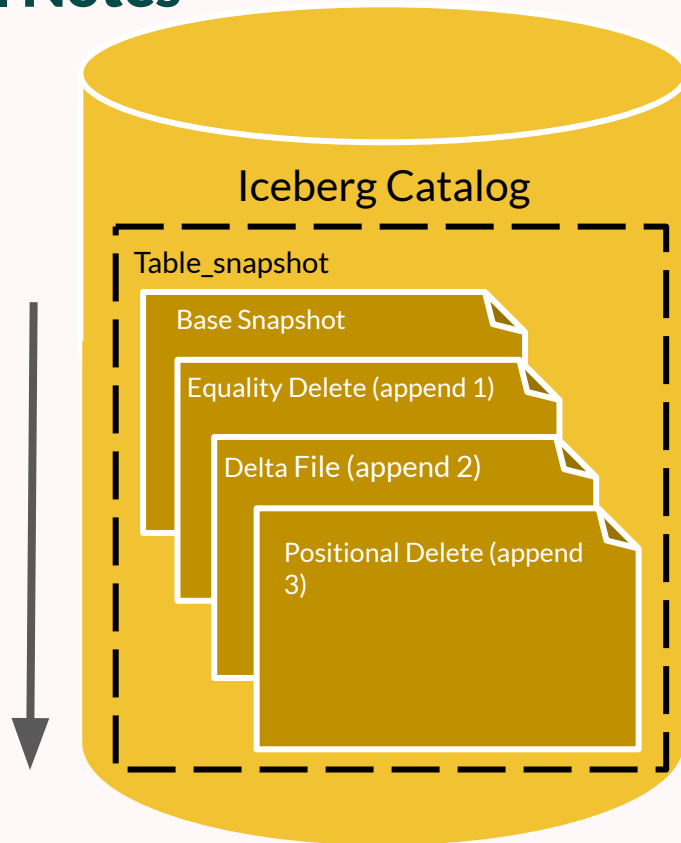


Desired Result		
id: 25	value: b	op: U
id: 30	value: alpha	op: U
id: 45	value: d	op: U

❖ **Duplicates within the datafile are filtered out**

Design : Iceberg V2 storage – Additional Notes

- Delete files effects performance
 - Positional Deletes → fast
 - Equality Deletes → slow
- Lots of small files
- **Regular maintenance** required to optimize the table
 - Rewrite data + delete files



Results

< 7min

Ingestion Time

3-40 min

Snapshot Compaction Time

Results

<7 min

Ingestion Time

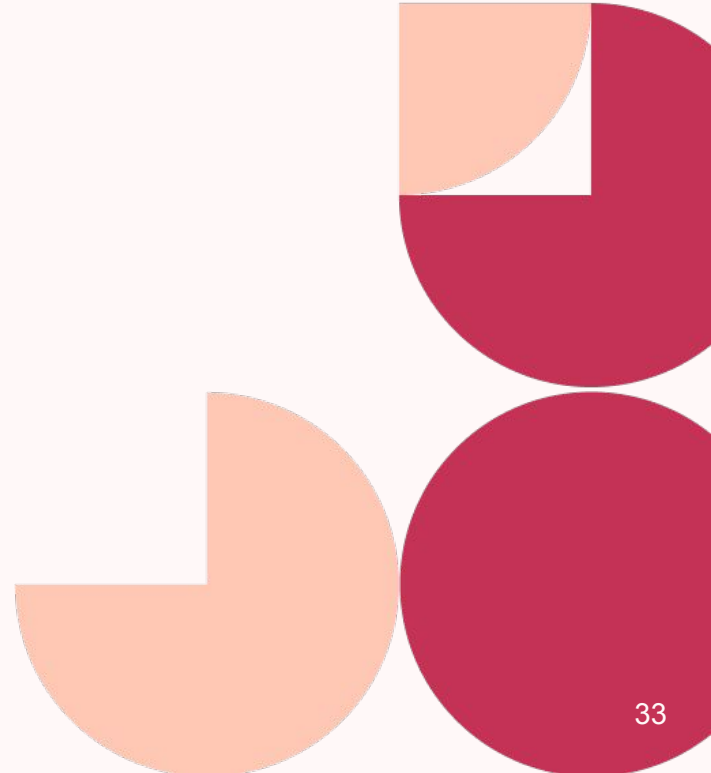
3-40 min

Snapshot Compaction Time

- **Lots more levers** (weigh cost vs performance)
 - Parallel partition rewrites
 - Max file group size (limit the input data to rewrite)
 - Machine size

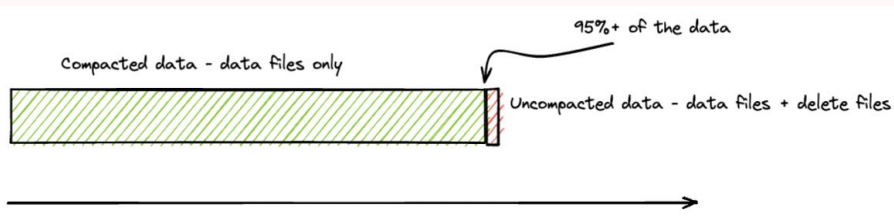
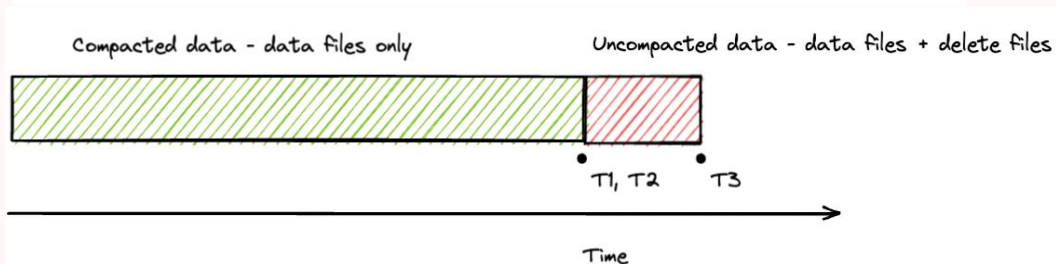
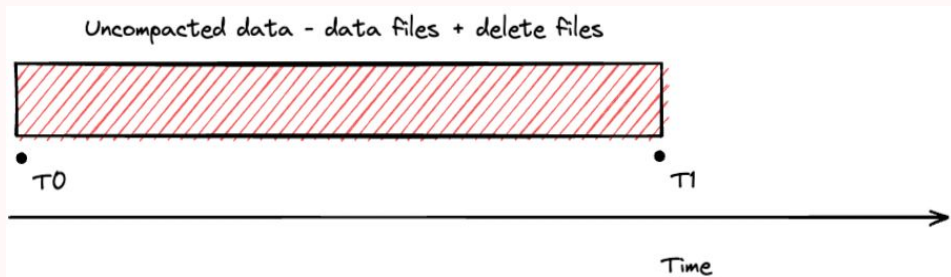
Goals

- ☑ Scalable ingestion
- ☑ Scalable **snapshot production**
- ☑ More accurate updates
- ☑ Delete capture
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Reflection

- Two ingestion modes
 - Initial full state dump
 - Steady state continuous ingestion
- Two categories of tables
 - **Fast** to query the **uncompacted**
 - **Slow** to query the **uncompacted**
- Performance cost of keeping delete files around
- Aggressive rewriting



Reflection

- Very happy with compaction times + performance
 - Having more levers is important
- Fast vs Slow uncompactd tables
 - Different tiers of tables
 - Hiding uncompactd tables behind a view or tag
- Iceberg's abstraction hides how data is represented on disk, which makes this possible
- Increased complexity of moving to streaming



Thank You!

